Representing Change by Aspect*

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Abstract. We propose the application of aspect-oriented programming to software configuration management. We believe it could improve the change control by providing a new basis for reasoning about a change. To demonstrate this, we designed an abstract-oriented extension to procedural languages where a change is represented by an aspect. Consequently, a change gains the properties of an aspect: it becomes well-localized and separated from the (unchanged) base program. This goes beyond the current capabilities of configuration management methods and tools: the aspect representing the change can be applied to other versions of the program (possibly to different programs).

Keywords: aspect-oriented programming, change control, change representation.

1 Introduction

Software systems are developed and evolved in a series of changes. Changes arise as requirements are extended, reformulated, dropped or corrected, as faults are discovered, and in many other situations. The change is often required also due to a need for adapting the product to the user’s context. We are witnesses of growing cooperation among software development companies. Many (often distributed) teams work on the same release of the software system in parallel. In such a situation, change control becomes even more important.

The level of change control support provided by the existing software configuration management tools varies significantly. Hence, if two companies decide to cooperate, there is a big chance that they would have different tools that provide different repository items representation, different structure representation, etc. The companies can also have different configuration management process established, including, for example, different branching and merging strategy, what even more complicates keeping track of changes in the source code.

We propose a solution to some of these problems by treating a change at the source code level and by expressing it explicitly. To achieve this, we employ the aspect-oriented programming, a new approach to programming aiming at separation of crosscutting concerns (see Section 3).

The rest of the paper is structured as follows. First, we put our approach in the context of the existing versioning models (Section 2) and the aspect-oriented programming (Section 3). Then we present an abstract aspect-oriented extension to procedural languages (Section 4). Subsequently, we show how this extension can be concretized to VBScript language (Section 5). Finally, we draw some conclusions and point some directions for the further work (Section 6).

2 Version Models and Change

A version model defines the entities to be versioned, version identification and organization, as well as operations for retrieving existing versions and constructing new versions [4]. Several version models are described in the literature and used in existing configuration management tools. We focus on the core issue of versioning, namely the organization of version space, or to be more specific—the version description and representation. Our main interest is to improve change control.

According to the entities being handled, the version models are classified into state-based and change-based. State-based models focus on the states of versioned items. In such approach, versions are usually described in terms of revisions and variants [2]. A configuration item (the smallest unit of a system taken under version control) is maintained usually at the file level. The change in state-
based models can be described as the difference between two versions. Many commercial systems are state-based (e.g., Microsoft Visual Source Safe, Rational ClearCase, PVCS) [3].

The problem with state-based models is that a change is maintained implicitly, during the modification of a branch. Merging can be viewed as a re-application of all the changes to the branches being merged. This requires the extraction of the changes from the branches and their subsequent application to the base.

In change-based models the change is treated as a first class entity and managed explicitly by a developer, either manually or by a tool. A version is considered as the result of the application of changes to a baseline. There are several commercial change-based software configuration management systems that have the ability to track the logical changes rather than individual file changes (e.g., Continus/CM, CCC/Harvest). They treat the change at the level of the source code lines or at the level of the file versions. Accordingly, they allow to create change sets or change packages [13]. A change set consists of the changed code lines. A change package contains references to the file versions that are the compositions of logical changes.

In a change-set model, changes are combined freely to construct new versions according to the requirements. In [4] such approach is denoted as change-based intensional versioning. The use of change packages is denoted as change-based extensional versioning, because version set is defined explicitly by enumerating its members. In this case, each version is described by changes relative to some baseline.

Another change-based approach is based on change identification by language constructs. This can be denoted as language-aware approach: the change is handled by directives for source code inserting, deleting and editing augmented with the attributes of the change (e.g., who and when made the change, etc.). An example of this approach is VTML (Versioned Text Markup Language) [9] or conditional compilation.

The conditional compilation enables to use the preprocessor directives to control the code fragments’ visibilities. In this case, all changes (fragments) are stored in one file, which is hard to maintain. Management of fragments’ visibilities is necessary for improving change control. This approach is used, for example, in the EPOS system [5].

The change-based systems are not so widely used as the state-based ones. The main reason is that developers think rather at the version-state level. However, nowadays many state-based systems are being extended to provide the change-based functionality (e.g., Rational ClearCase) [8]. The objective is to improve change management and traceability of the change request in a software development process.

The change representation influences the change control procedure, which consists of the four major steps [1]: checking whether the change is needed, analysis of causes that led to change, planning the change, and change implementation. In the context of the change control procedure, we are concerned with the change implementation.

The problem with the surveyed change-based approaches is the granularity of the logical change. As we mentioned, some of them treat the logical change as the individual lines of the source code, while other are based on representing change by preprocessor directives.

3 Aspect-Oriented Extensions

The main idea of the aspect-oriented programming (AOP)—separation of concerns by separating the cross-cutting concerns called aspects from the basic functionality crosscut by them—is carried throughout several independently developed approaches [10, 12]. Among them, Xerox PARC AOP [14] holds a significant position. Further in the text by the AOP we mean actually the Xerox PARC AOP.

AOP appeared as a reaction to the problem known from the generalized procedure languages [7], i.e. programming languages that use the concept of the procedure to capture the functionality. In such languages the program code fragments that implement a clearly separable aspect of a system (such as synchronization) are scattered and repeated throughout the overall program code that, in advance, becomes tangled. AOP aims at factoring out such aspects into separate program units called by the same name: aspects. Aspects crosscut the base code in places called join points. These must be specified so aspects could be woven into the base code by the program called weaver.

The join points can be static or dynamic. Static join points can be identified in the program text itself. They can be specified in terms of a programming language syntax alone. An example of such a join point is the beginning or end of a method or procedure body. Dynamic join points are available at run time only. For example, a method reception by an object is a dynamic join point. In the weaving process, the static join points are resolved by a simple program code insertion, while dynamic join points can be resolved at run time only.

The special language constructs used to capture the aspects and join points are known as the aspect-oriented extension of the base language. The two types of aspect-oriented extension regarding its relationship to the base

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1It is possible that some workaround for the existing activity could be more effective than the change itself.

2Besides the procedural languages, these include functional and object-oriented languages as well.
language can be distinguished: homogenous and hetero-
geneous. The homogenous extension, besides for some
additional constructs, relies on the base language to the
greatest possible extent, while the heterogeneous exten-
sion introduces a whole new language for capturing the
aspect-oriented part of the program. In general, there can
be several independent aspect-oriented extensions, han-
dled by the same or by separate weavers.

Not unlike programming languages in general, an
aspect-oriented extension (including the corresponding
weaver) can be designed to solve a specific problem, such
as the one presented in [7] (the filtering example), or to
serve a general-purpose, as the AspectJ language [15],
which is a homogenous, general-purpose aspect-oriented
extension to Java. While aspect-oriented extensions pro-
vide a new way of programming, they do so only in the
context of the language they extend. In other words, AOP
is a multi-paradigm approach in its very nature [12], and
AspectJ can be viewed as a multi-paradigm language [11].

4 Aspect-Oriented Extension for
Change Representation

As it was discussed in Section 2, current configu-
ration management approaches do not offer a satis-
factory change representation regarding the change maintenance
and re-applicability to different branches. The use of
AOP enables to maintain changes explicitly by captur-
ing a change into an aspect.

In order to enable change representation by aspect, the
aspect-oriented extension to a given programming lan-
guage should be provided. Since the changes are actually
changes of the program text, all the join points will be static.
Further, the aspect-oriented extension should be
homogenous—to preserve the base language constructs,
and general-purpose—to cover all the types of changes
(which depend on the base language). Also, the join point
description should not affect the base program.

To illustrate aspect-oriented approach to change rep-
resentation, we developed an aspect-oriented extension
(inspired by AspectJ) to procedural languages. Proposed
language constructs are presented in Fig. 1. Different
type styles are used to distinguish among the keywords,
required parts and optional parts.

The aspects are placed into modules, possibly together
with the ordinary procedures which can be called from
within the aspects, i.e. inside of the block parts. The
block parts must be parsed either by the weaver, or by
the original language parser.

The introductions are used to introduce new procedures
and variables into modules (M_i). The advices enable per-
forming a command block before, after or in place of the
procedures determined by a specified set of join points,
so-called pointcut. While before and after advices are
simple, the around advice requires some explanation. It
enables to run an initial block block, then to proceed
with the next action, which is either another aspect, or
the original procedure body, in case there is no other as-
pect affecting the procedure. The optional return_clause
in after and around advices enables to modify the return
value (if the procedure returns one) before it is actually
returned to the caller.

The pointcut specification is built out of the point-
cut primitives (listed in the bottom of Fig. 1) using
the logical operators and and or. The parentheses
can be used to declare the priority of subexpressions
evaluation. The first two primitive pointcuts, module
and withincode, designate all the join points within
the modules M_i and procedures specified by the pro-
cedure signature, respectively. The calls pointcut
primitive designates all the procedure calls specified by
the procedure signature. The definitions pointcut
primitive designates the actual definitions, i.e. bodies of
the procedures specified by the procedure signature
(see Fig. 5 for an example). A before advice to a de-
definitions pointcut will insert its code after all the
declarations of variables in the specified procedure(s)
placed before the first non-declaration statement.

The wild cards * and .. can be used in procedure
signature to denote any string of characters and
omitted arguments, respectively. This convention is used
in AspectJ, e.g. * p*(int, *) denotes all the methods
whose name starts with p, with one int argument and
one argument of any type, returning a value of any type.
The most general signature—denoting all the methods—
is then * * (*..*).

Up to now we said nothing about the optional argu-
ment_list in advices. It is used to access the arguments
of the procedures denoted by the pointcut. Suppose we
want to make a before advice to the following C function:

\[
\text{int } f(\text{int } i) \{ \text{return } i*i; \}
\]

Consider these two advices:

1. \textbf{before()}: \textbf{definitions (int } f(\text{int)) \{ i = i + 1; \}}

2. \textbf{before(int } x): \textbf{definitions (int } f(x))
   \{ x = x + 1; \}

Both advices seem to do the same thing; they add one
to f's argument before proceeding with the rest of f's
body.\(^4\) However, if we rename i in function f to j, the
first advice will fail to satisfy our intention (moreover,
it will produce a syntax error), so the second version is
obviously more robust.

\(^3\)Since pointcuts are the sets of join points, the and and or opera-
tors have the meaning of set intersection and union, respectively.

\(^4\)This is different from AspectJ where the method body is not
visible to advices.
Introductions:

introduction M_1,...,M_n \{ block \}

Advices:

before\((arguement_list)\): pointcut \{ block \}
after\((arguement_list)\): pointcut \{ block return_clause \}
around\((arguement_list)\): pointcut \{ block, proceed... block, return_clause \}

Pointcuts:

pointcut pointcutName \((arguement_list)\): pointcutSpecification

Pointcut primitives:

modules(M_1,...,M_n)
withinCode(p\(rocure\_\text{signature})
calls(p\(rocure\_\text{signature})
definitions(p\(rocure\_\text{signature})

Figure 1: Aspect-oriented extension to procedural languages.

The proposed aspect-oriented extension is capable of describing the following types of changes:

- introduction of a new procedure or (global) variable into the module;
- extension of a procedure by a code before, after, or instead of it;
- change to the procedure arguments and return value.

What all of these changes have in common is that they are all about the functionality. The changes that cannot be described at the level of functionality are very hard (or impossible) to deal with using the aspect-oriented approach. These include renaming a procedure or variable, adding a white space or comment, changing the position of a procedure in the source code, etc.

A version is obtained by weaving the aspects that capture the change into the base program. Since this version might become a subject of modification as well, it should be human readable. This is different from the AOP itself where the process of weaving yields only an intermediate product not intended to be read by a human. The aspect-oriented extension proposed principally satisfies this requirement, since it relies on static join points only.

5 Case Study: Script Customization

We will show now how the approach we proposed in the preceding section could help in solving a problem of synchronizing the local customization with the global version of a program in script languages by the means of an example. We will use VBScript-like syntax since VBScript is widely used as a language for dynamical content generation and design of the web pages. It is the core language for Microsoft’s asp pages. Further, some software houses use VBScript as the language for customizing their products, e.g. InteractCommerce corp.’s SalesLogix.

Suppose that two teams work on one system. The teams received change requests regarding the same script, shown in Fig. 2, which is a part of the system, at the same time, i.e. before synchronization of branches, as depicted in Fig. 3. The purpose of the script is to extract the list of sales opportunities from the opportunity table in the test database. The change request received by the first team is about extending the list of opportunities by the list of products. A new recordset, as well as the SQL statement and several lines of VBScript code must be added in order to accomplish the task of extracting the records from the table and generating the page containing the data.

The modified script is presented in Fig. 4; some commands the same as in Fig. 2 have been omitted (indicated by ellipsis). The code between the change and end change comments can be separated into the aspect module, as presented in Fig. 5. The affected module is specified by the modules designator in both advices. The declarations of additional variables are provided in the before advice. The conjunction of the definitions and modules designator states that the sequence of variable declarations in the advice is to be merged after all the declarations in the main procedure which are placed before the first non-declaration statement. The sequence of the directives to be run after the ro.close method in-
vocation (specified by the \texttt{calls} designator) within the original \texttt{main} procedure (specified by the \texttt{withincode} designator) has been enclosed into the \texttt{after} advice. The result of the merging, i.e. weaving, will be the same code as displayed in Fig. 4.

The other change request addressed to the second team resulted in the script shown in Fig. 6. The second team’s change consists of adding the \texttt{while} loop for updating the \texttt{applied} and the \texttt{date_of_application} fields for each record in the opportunity table, and of adding the sequence of commands that generate the list of marketing campaigns from the \texttt{marketing} table. We can apply the aspect from Fig. 5 to the code in Fig. 6 without any change.

However, that change could have been separated into the aspect, too, if both teams used the aspect-oriented approach. In that case, we would simply apply the two aspects subsequently in order to obtain both functionalities.\footnote{The priority of aspects is not significant here, but this is not so in general.}

\section{Conclusions}

We proposed a new approach to change-oriented versioning based on the aspect-oriented programming. The contribution of this paper is the proposal of the technique aimed to simplify change control by reifying the changes into language-level entities: a change is represented by an aspect and maintained explicitly by a developer.

A homogenous, general aspect-oriented extension has to be provided for a given programming language first. For the purposes of our approach, it is sufficient if this extension supports static join points. Since procedural, functional and object-oriented languages are easily extended to support the AOP with static join points, this approach is low-cost. We proposed such an extension to procedural languages. Moreover, it can be expected that general aspect-oriented extensions to other programming languages will be developed and provided for the sake of the AOP itself, so no additional effort would be necessary to employ this approach in such languages. This can be denoted as \textit{self-supported change management}: a change is represented by the constructs that are a part of the programming language itself.

We assume this as one of the main advantages of the proposed version space representation. It provides a new base level for the change control; it is a move from the change control at the line level to the one at the programming language semantics level. In small software projects it is directly usable even without a software configuration management tool. The change comprehension and ori-
before(): modules(script) & definitions(main)
    begin
    Dim rp 'Recordset object
    Dim s2 'Select statement
    Dim c 'Command object
    end before

after(): modules(script) & withincode(main)
    & calls(ro.close)
    begin
    s2 = "SELECT * FROM opportunity_product"
    Set c = Server.CreateObject("ADODB.Command")
    c.ActiveConnection = con
    c.CommandText = s2
    Set rp = c.Execute
    call gener_data(rp)
    rp.close
    c = Nothing
    end after

Figure 5: The change separated into the aspect.

tentation in the source code is easier because the change
is well-localized in the aspect and need not be searched
for. A change is possibly re-applicable as is or with some
adaptation of the aspect involved (white-box reuse). Ac-
tually, the aspect can be applied to a completely different
module than it was intended for by a simple modification
of the pointcut.

Aspect-oriented approach can be used also in post-deployment configuration management [6] for
parametrization (modification of a software system to
take into account the local site context). The local con-
text can be represented by an aspect. The application
of the relevant aspects provides customization of the new
product version according to the local context (developed
for the previous version). Obviously, new aspects will be
also created, in order to customize new features in the
current version of the product.

Our approach can be used with existing software con-
figuration management tools. Moreover, our approach is
independent of the model employed by software configu-
ration management tools. An aspect is a separate item,
so it can be handled in both basic version models (state-
based and change-based). It also supports the implicit
long transaction maintenance because aspect itself rep-
resents a change and it is up to the developer to decide when
the change should be committed. As aspects can be sim-
ply plugged in or out before the compilation, adding of an
individual change into a version or substracting a change
from a version (similarly as in the change-set approaches)
is simple. Even more, the aspects can be combined into
change packages. A change request could be then directly
assigned to the corresponding aspect or change package
(indicated by the appropriate identifiers).

Our work is now oriented toward a deeper elaboration
of practical use of the proposed approach. Some ad-
ditional mechanisms should be added to manipulate ver-
sion history and versions themselves. In order to be able
to control the changes effectively, some meta-data should
be stored within each change (e.g., who and when made
the change). In order to follow a version history, the
meta-data related to changes should be processed and in-
terpreted.

On the other hand, we are already able to partially
track the version history, but it is difficult to determine
which version was created by a developer and which is
just a potential version when storing only changes. The
potential versions can be obtained by applying the combi-
nations of the aspects. However, not all potential versions
make sense [4].

An additional problem is that the aspect representing a
change can become a subject of change, too. As a con-
sequence, a method of dealing with the change of a change
should be proposed. This problem arises in any change-
based version model, of course, but a special method is
needed here because our approach works at other level
of change control than traditional change-based version
models.

sub main
    . . .
    '***change of declarations***
    Dim rm 'Recordset object
    Dim com 'Command object
    Dim str 'String - select statement
    '***end of change declarations***
    Set con = . . .
    . . .
    call gener_data(ro)
    '***change***
    While Not ro.EOF
        ro.Fields("applied") = 'T'
        ro.Fields("date_of_application") = Now
    Loop
    str = "SELECT * FROM marketing"
    Set com = Server.CreateObject("ADODB.Command")
    com.ActiveConnection = con
    com.CommandText = str
    Set rm = com.Execute
    call gener_data(rm)
    rm.close
    rm = Nothing
    com = Nothing
    '***end of change ***
    ro.Close
    . . .
    end sub

Figure 6: The change performed by the second team.
References


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